

CS118: Computer Network Fundamentals

Lecture-1: introduction

CS118: explains (roughly) how the Internet works

- ◆ Internet: a huge, complex network of networks
- ◆ Divide-and-conquer
 - Figure out how many major parts,
 - Learn one piece at a time
- ◆ Your job:
 - Read textbook, think, collect a list of questions
 - review every lecture slide deck after each class
 - Ask questions in class/office hours/via Piazza
 - Practice what you learn through homework and projects

James F. Kurose | Keith W. Ross

COMPUTER NETWORKING

A TOP-DOWN APPROACH

Eighth Edition



Brief Contents

Chapter 1	Computer Networks and the Internet	1
Chapter 2	Application Layer	81
Chapter 3	Transport Layer	181
Chapter 4	The Network Layer: Data Plane	303
Chapter 5	The Network Layer: Control Plane	377
Chapter 6	The Link Layer and LANs	449
Chapter 7	Wireless and Mobile Networks	531
Chapter 8	Security in Computer Networks	607
	References	691
	Index	731

xix

Course assignment and due schedule

Midterm	In-class, Wednesday Feb 5 (Location TBD)	
Final	3:00PM-6:00PM Saturday March 21 (Location TBD)	
Homework	Release: on Thursday of week 1, 3, 5, 7; Due: 11:59pm Tuesday of week 3, 5, 7, 9.	
Project 0	Release: Monday Jan 6, 2024 (Week 1) Due: 11:59pm Wednesday, Jan 15, 2024 (Week 2)	1.5 weeks
Project 1	Release: Thursday Jan 16, 2024 (Week 2) Due: 11:59pm Sunday, Feb 16, 2024 (Week 6) Grading: auto-grading script (sample tests will be provided to let everyone test their code before submission)	4.5 weeks
Project 2	Release: Monday, Feb 17, 2024 (Week 7) Due: 11:59pm Sunday, March 16, 2024 (Week 10) Grading: auto-grading script (sample tests will be provided to let everyone test their code before submission)	4 weeks

FOR ALL OTHER COURSE INFO, PLEASE SEE [HTTPS://BRUINLEARN.UCLA.EDU/](https://bruinlearn.ucla.edu/)

Course workload and grading

- ◆ Bi-weekly homework assignments
- ◆ 3 programming projects,
 0. UDP socket (individual)
 1. Reliable data delivery (2-3 people team)
 2. Secured reliable data delivery (2-3 people team)
- ◆ Midterm and final exams (cheat sheets allowed, 2 pages double sided)
- ◆ **Strict Grading Policy**
 - Homework: do it yourself; no credit for late submission
 - Project: 20% credit reduction per late day
 - *No make-up exam*

Homework	20%
Programming Projects	25% (5/ 10/ 10)
Midterm	25%
Final exam	30%

2% extra credits based on piazza

1% extra credits course evaluation and TA/LA feedbacks

Class Policy

The following actions are **strictly prohibited**

- ◆ Posting/sharing/selling class material, with or without answers, to anyone outside this class, during or after this quarter.
- ◆ Use of old homework/midterm/finals in doing homework or exams
- ◆ Use ChatGPT in doing assignments
- ◆ Making your project code publicly available either during or *after* this quarter
 - *you must use private repository* on GitHub or GitLab

Hints for Getting Good Grade

- ◆ Review previous lecture slides
- ◆ Read textbook before coming to each lecture
- ◆ *Ask questions*
- ◆ Get your work done early
 - Lecture slides uploaded to BruinLearn *one day before the lecture*
 - Get HWs and projects done **before** the deadline

Let's get started

Today we cover the basic concepts in
Chapter 1 of the textbook

What is a Computer Network




What is a Computer Network





Coming to you



 **Restaurant**
★★★★★ 25m

 **Grocery Store**
★★★★★ 30m

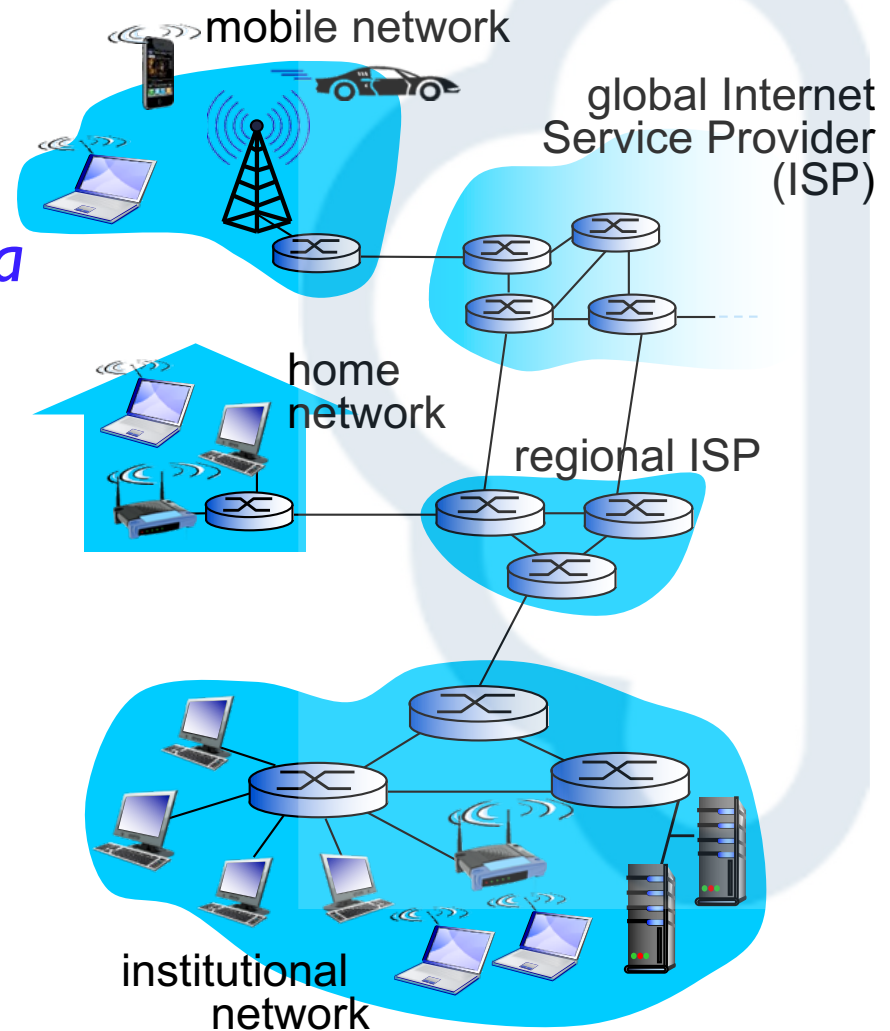
 **Hotel**
★★★★★ 40m

 **Transit**
1:30min 120m

Terminology



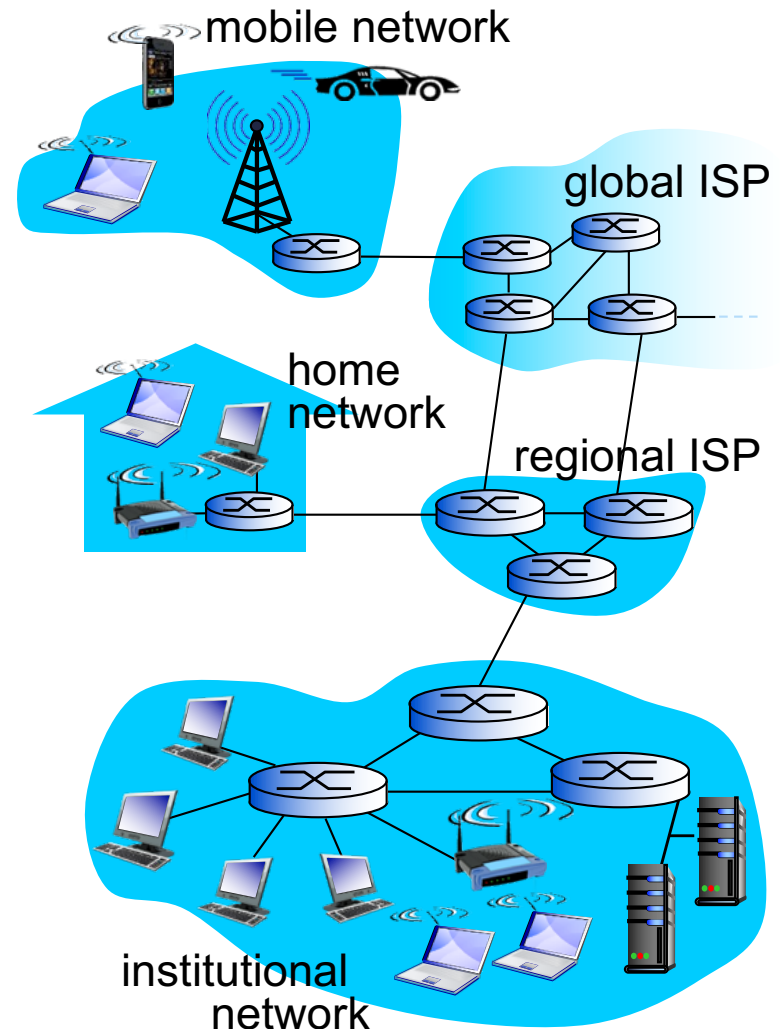
- billions of connected computing devices:
 - hosts* = *end systems*
 - running *network apps*
 - Apps send/receive *data packets*
- Routers* = *packet switches* inside network
- communication links*
 - fiber, copper, radio, satellite
 - transmission rate = *bandwidth (BW)*



Recent years witnessed rapid growth of giant cloud service providers

“Nuts and Bolts”

- ◆ *Internet*: “network of networks”
 - Interconnected ISPs, enterprise networks, now also cloud service providers
- ◆ *Protocols*: define how to send, receive packets
 - e.g., HTTP, TCP, IP, 802.11
- ◆ *Internet protocol standards*
 - RFCs: “Request for Comments”
 - <https://www.rfc-editor.org/rfc-index.html>
 - Developed by Internet Engineering Task Force (IETF)
 - IEEE Standards
 - W3C (World Wide Web Consortium), and others



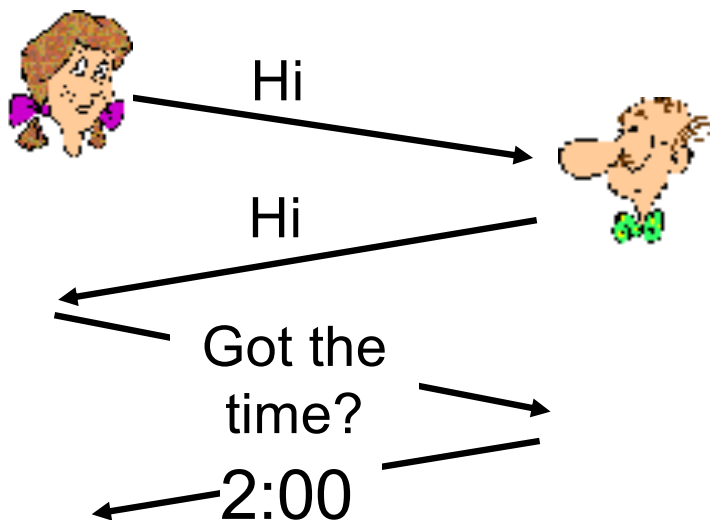
What is a protocol?

Traffic light protocol

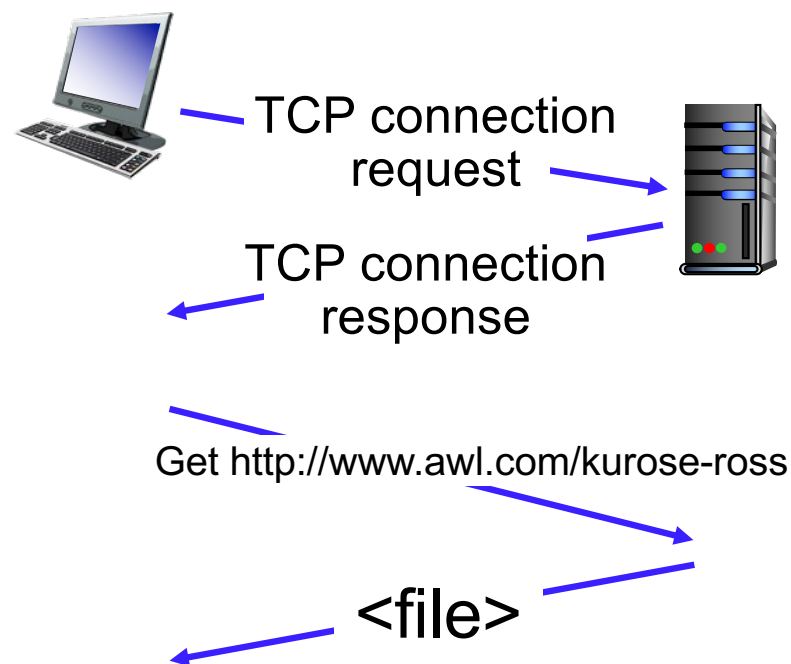
- ◆ Green: go
- ◆ Red: stop
- ◆ Yellow: slow down - stop

- ... specific messages sent
- ... specific actions taken when the messages received

human protocols:

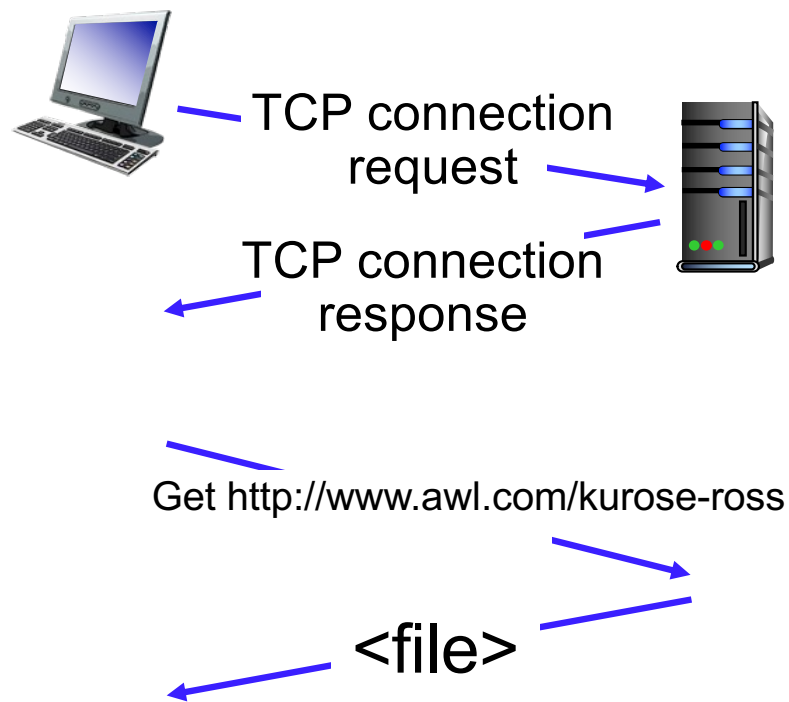


computer protocols:



Internet protocols

computer protocols:



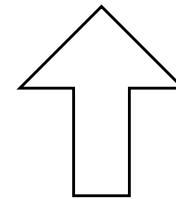
- ◆ Communication between machines rather than humans
- ◆ all communication activity governed by protocols

protocols define *format*, *order* of *packets sent and received* among network entities, and *actions taken* on packet transmission, receipt

Delivering data over the global Internet is a complicated process, involving many many steps

How to get the work done: divide and conquer

Group functions to a few modules



How many?

Internet protocol stack

◆ Application layer protocols

- Support data exchange between application processes
- Example: SMTP, HTTP, DNS
(Simple Mail Transfer Protocol)

◆ Transport layer protocols

- handling delivery reliability, multiplex within a host
- Example: TCP, UDP

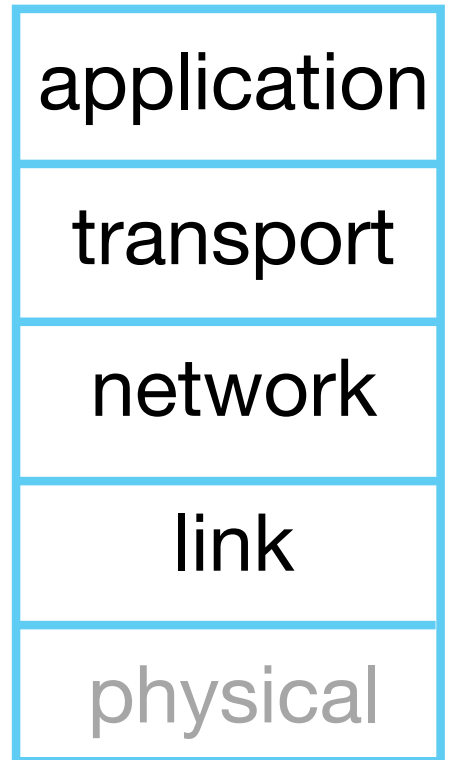
◆ Network layer protocols

- forward packets from source to destination
- Example: IP

◆ Link layer protocols

- transfer data between directly connected network elements
- Example: Ethernet protocol, WiFi

◆ Physical layer: bits “on the wire”



Application View

apps
trans
net
link
phy

My Laptop -
Running web
browser

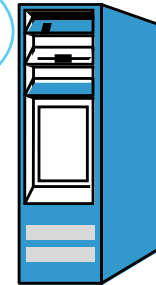


client

(Chrome, Safari, Firefox, ...)

Internet

Web Server
www.cnn.com



server

(Apache, GWS, ...)

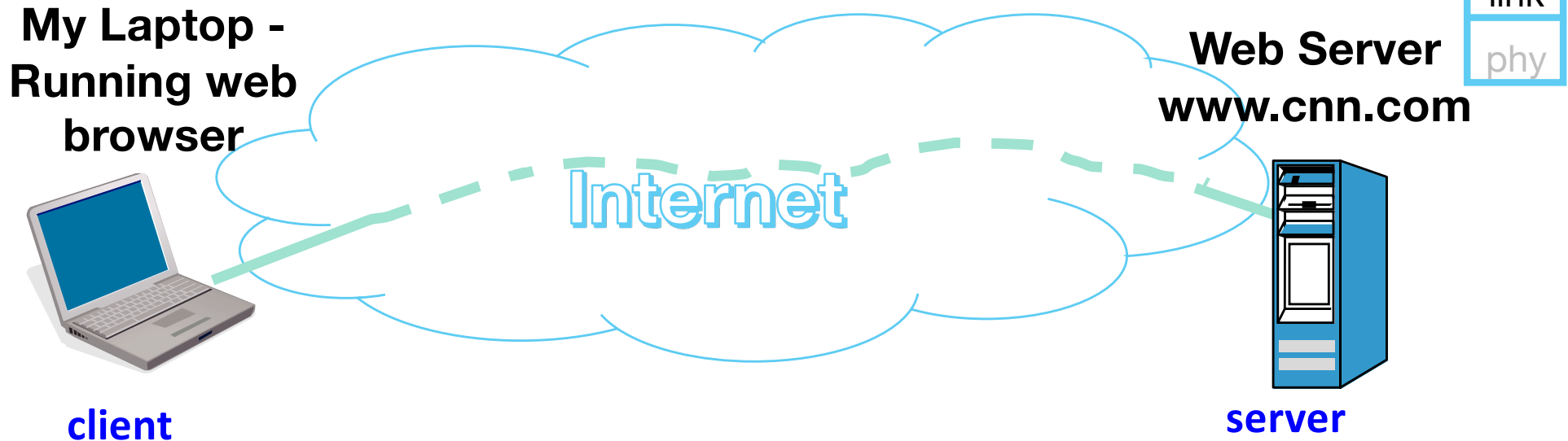
These are **application programs**

They talk to each other using **application protocols** (web protocol: HTTP)

Application protocols

- Assume network can send data to any hosts on the Internet
- Don't know/care how data is sent, and assume all data delivered reliably
- Runs on top of a transport protocol

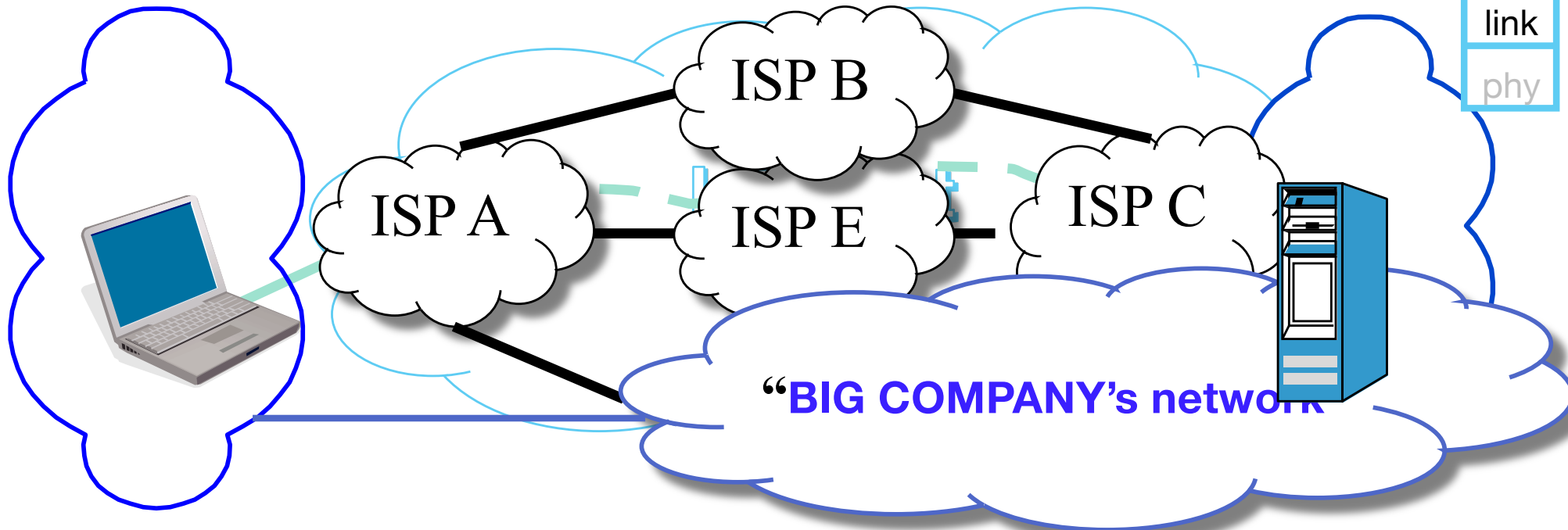
Transport View



- ◆ **A transport protocol's job: delivering data** between the two communicating ends
 - *Don't know or care about which paths data may traverse through the network*
- ◆ Multiple transport protocols exist, each offers somewhat different functions (e.g. reliability, congestion control)

Actually, transport protocols don't do delivery → network protocol's job

Network Layer View



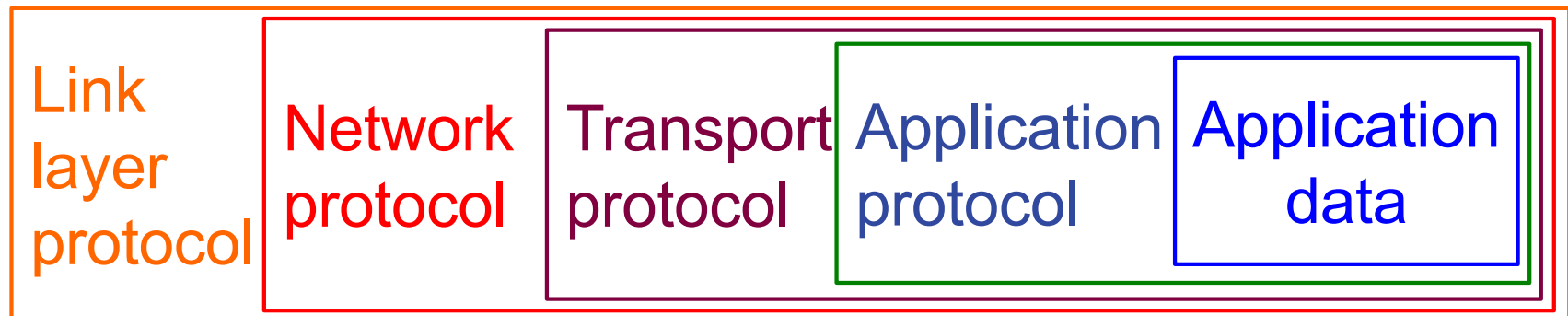
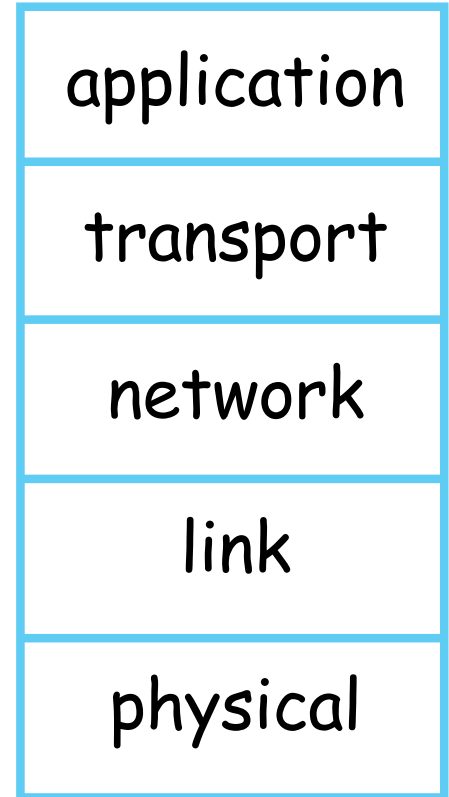
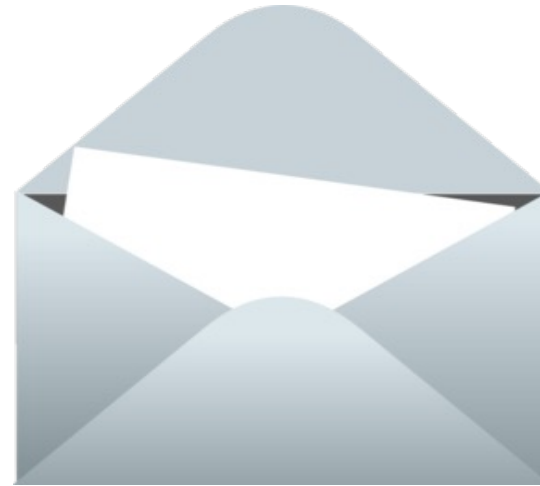
- ◆ **network protocol's** job: forward packets from source to destination host
- ◆ A really hard problem: the Internet is large, run by many different parties
 - connection from laptop to CNN.com:
WiFi → campus backbone → local ISP → other ISP → CNN website

Link Layer View



- ◆ **Link layer's job:** Get a packet transmitted across some communication medium to **next hop**
- ◆ Different medium → different link layer protocol

What protocol “layer” really means

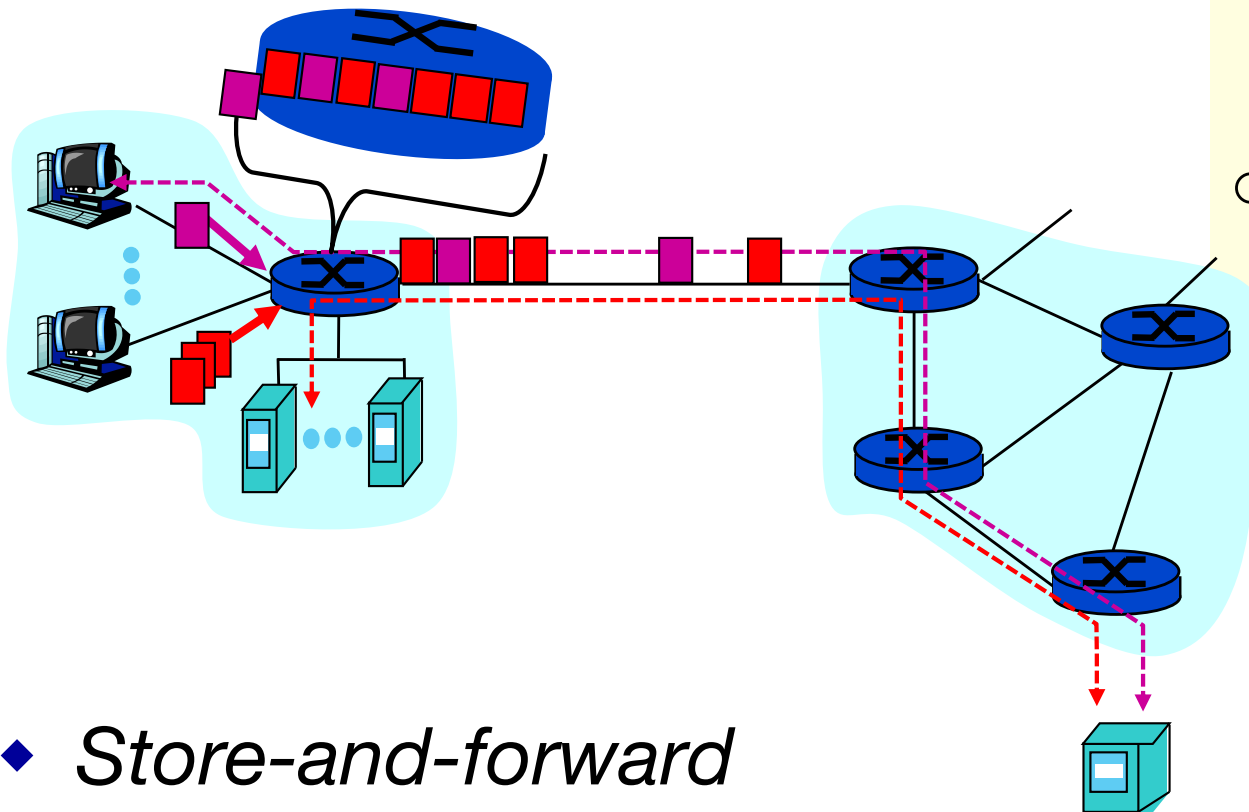


(Tentative) Schedule of the Quarter

Week:		1	2	3	4	5			
Mon		1/6 Course intro BW& delay	1/13 HTTP	1/20 Martin Luther King Jr. Day	1/27 Transport protocols	2/3 Congestion Control			
	Wed	1/8 Socket programming, Web & HTTP	1/15 DNS	1/22 DNS	1/29 TCP	2/5 Midterm			
		6	7	8	9	10			
Mon		2/10 Security 101	2/17 Presidents' Day	2/24 Routing algorithms & protocols	3/3 Routing in the Internet	3/10 Hubs and switches			
	Wed	2/12 Internet Protocol (IP)	2/19 Addressing, NAT, IPv6	2/26 Routing algorithms & protocols	3/5 Link layer (Ethernet)	3/12 Course review	3/21: Final Exam		

- The big yellow numbers indicate the chapter numbers in the textbook.

Packet Switching: *Statistical Multiplexing*



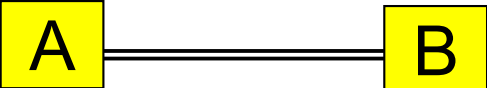
- Each node sends packets as soon as link available
- Receiver gets a full packet first, then forwards it towards the destination

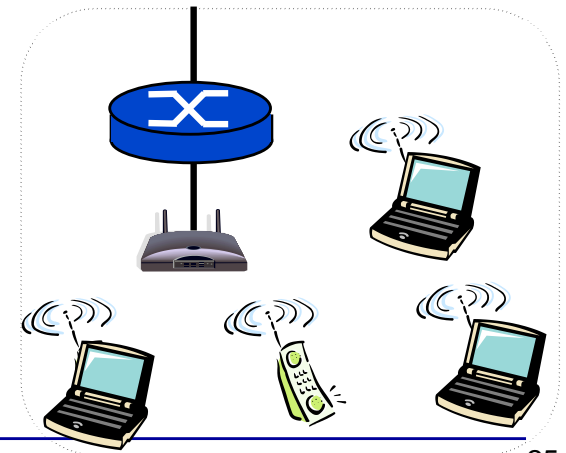
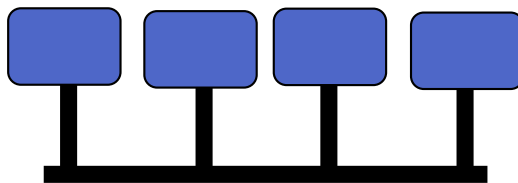
- ◆ *Store-and-forward*
- ◆ Packet switch can temporarily buffer up packets
 - Introduce *delay*
 - Packets get *dropped* when the queue is full

Network Performance

- ◆ 3 basic measurements
 - Throughput (bits/sec, Kbps=1000 bits/sec, Mbps)
 - Loss rate (% of packets lost)
 - Delay (sec, msec)

Throughput

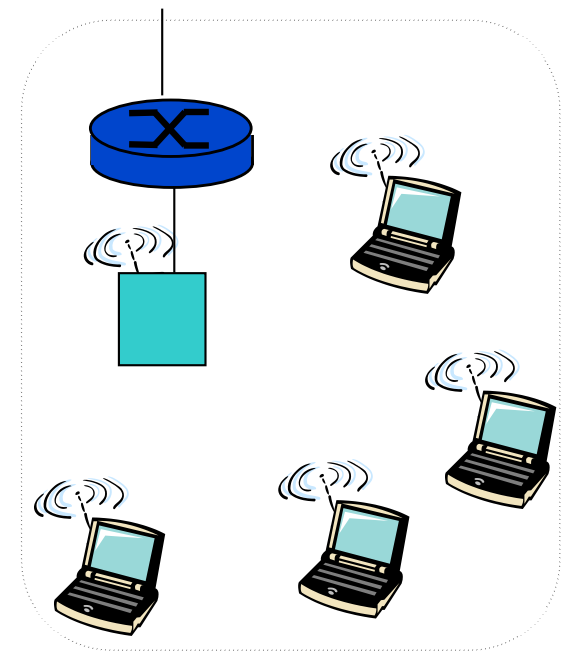
- over a single link: point-to-point 
 - Pumping data into the pipe: throughput = link bandwidth
- Multi-access:
a lot more difficult to measure, why?



Packet Losses

- ◆ Wired links
 - Loss due to transmission errors
 - Loss due to congestion
- ◆ Wireless links
 - Limited transmission rate
 - Higher (than wire) bit error rate
 - Host mobility: high variance in the number of hosts sharing the same wireless channel

Do users know there are packet losses?
Do users' performance get affected by losses?



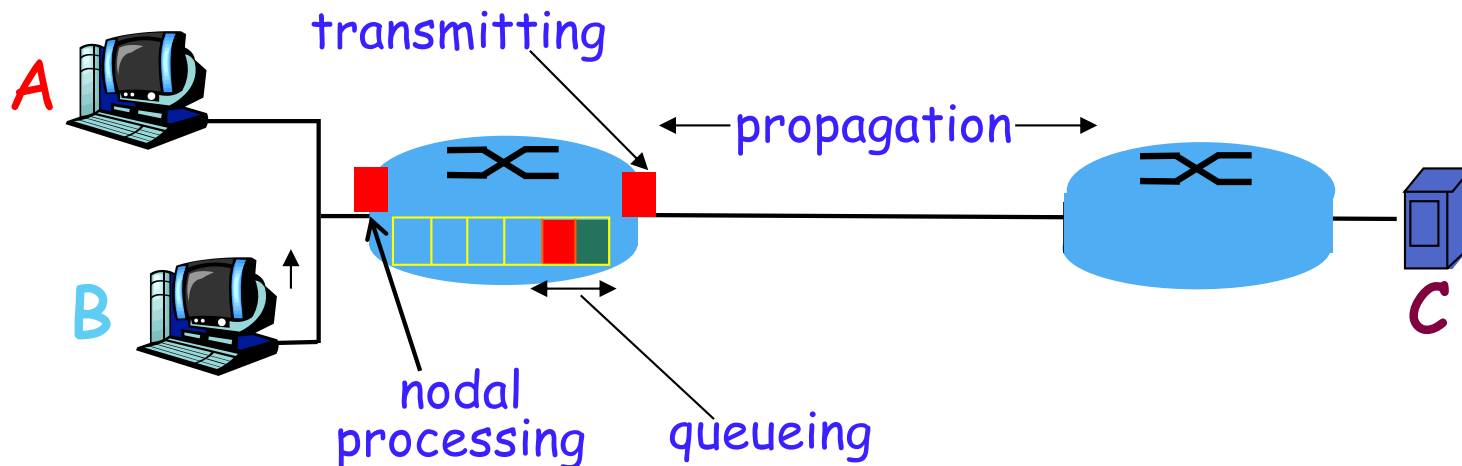
Delay in packet-switched networks

4 sources of delay at each hop


- ◆ node processing:
 - check bit errors
 - determine output link
- ◆ Queuing = #packets in queue
X transmission time
of each packet

Transmission = Length / rate
 R = link bandwidth (bps)
 L = packet length (bits)

Propagation = distance/sec
 d = length of physical link
 s = propagation speed in medium ($\sim 2 \times 10^8$ m/sec)

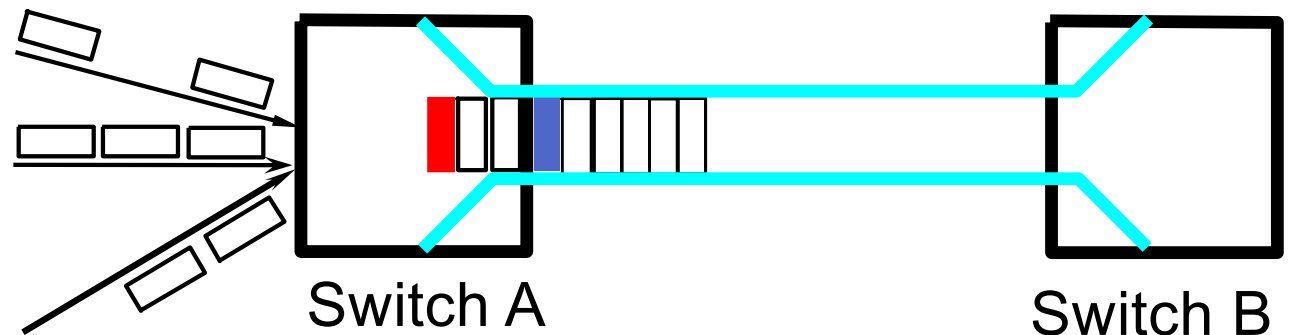


Example: calculating one hop delay

total delay (A  B) = ?

- ❖ Queuing delay = ?
- ❖ transmission delay = ?
- ❖ Propagation delay = ?

link length = 100 km
Bandwidth = 1 Mbps
packet size = 1000 bits
(all pkts equal length)



(2.0×10^8 meters/sec in a fiber)

Example: calculating one hop delay

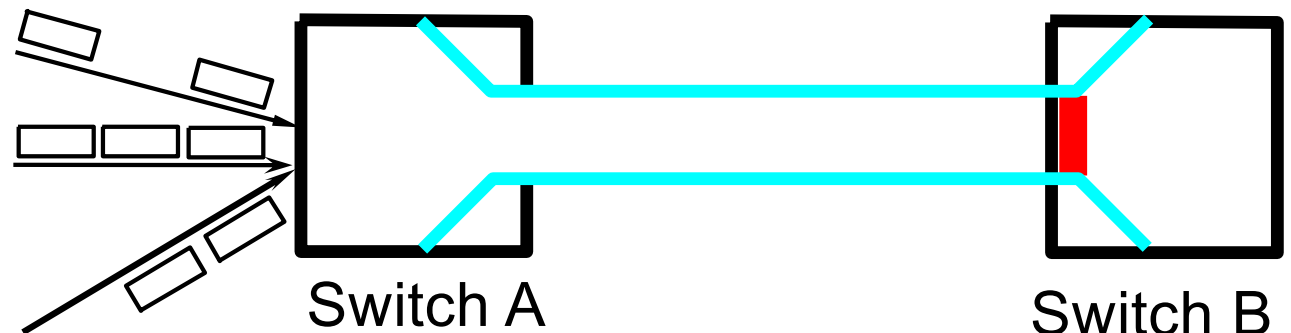
$$\text{total delay (A} \rightarrow \text{B)} = 1\text{ms} \times 2 + 1\text{ms} + 0.5\text{ms} = 3.5\text{ms}$$

❖ Queuing delay = **Waiting time for 2 pkts**

$$\text{transmission delay} = \frac{1000\text{bits}}{10000000\text{bits/sec}} = \mathbf{1\text{ msec}}$$

$$\text{Propagation delay} = \frac{100,000\text{m}}{2 \times 10^8\text{ m/sec}} = \mathbf{0.5\text{ msec}}$$

link length = 100 km
Bandwidth = 1 Mbps
packet size = 1000 bits
(all pkts equal length)



(2.0×10^8 meters/sec in a fiber)

Transmission vs. propagation delay

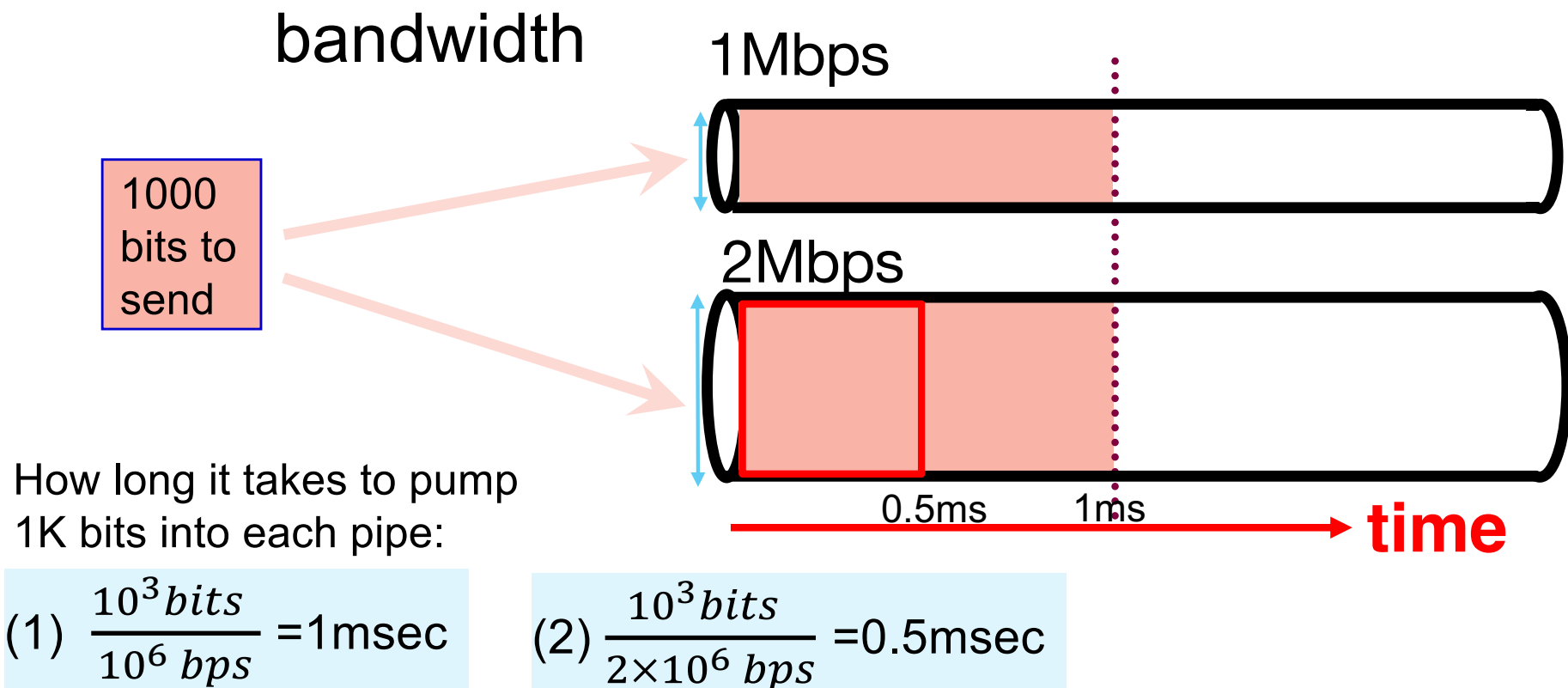
Transmission delay: L / R

R = link bandwidth (bit-per-second, bps)

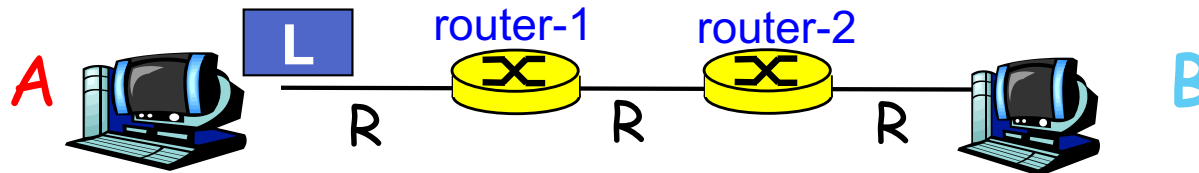
L = packet length (bits)

Propagation: d / s

d = length of a physical link
 s = signal's propagation speed in the medium ($\sim 2 \times 10^8$ meter/sec)



Packet-switching: store-and-forward

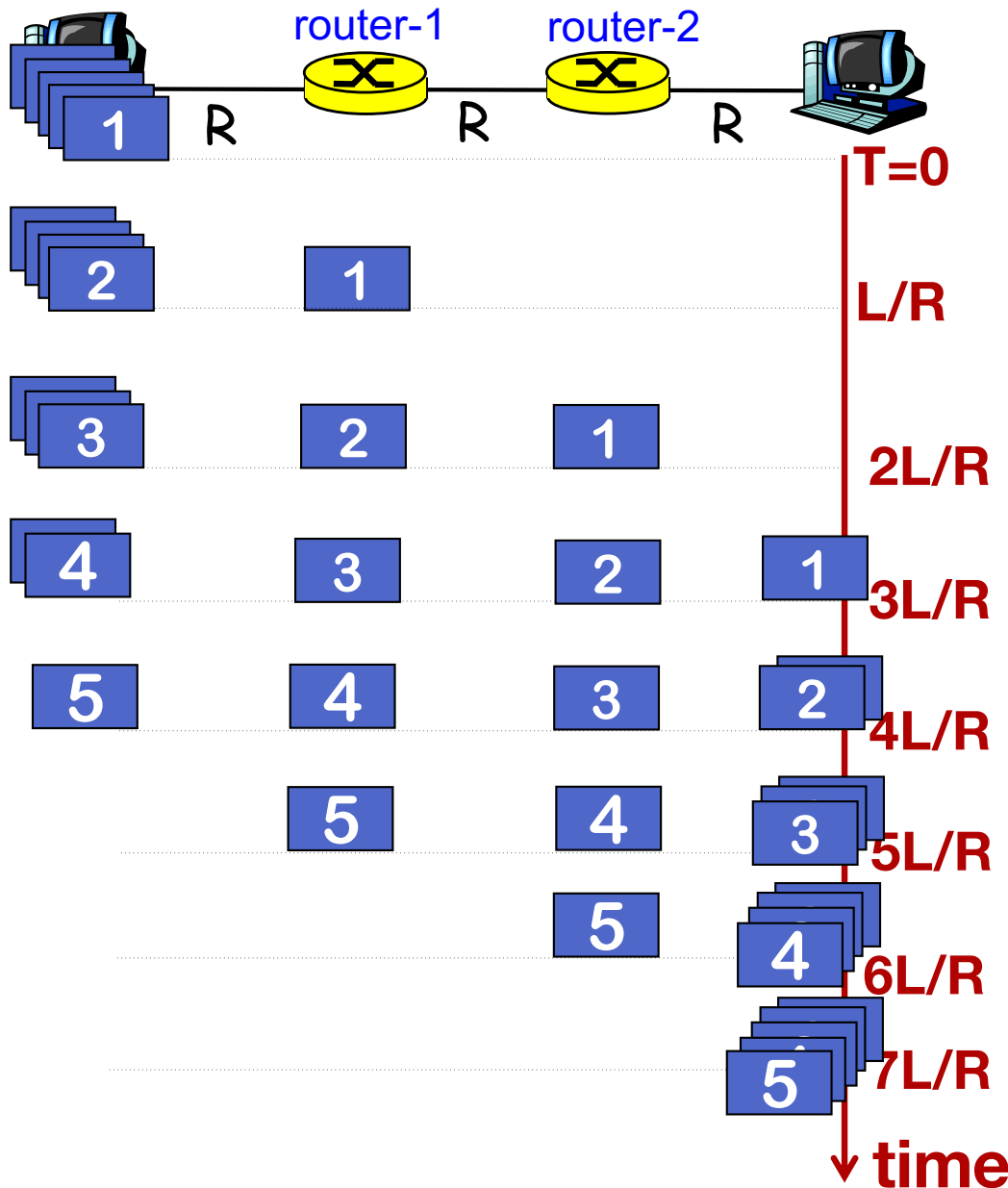


- ◆ Takes L/R seconds to transmit (push out) packet of L bits on to link of R bps
- ◆ Entire packet must arrive at router before it can be transmitted on next link: *store and forward*

Example 1: send L $A \rightarrow B$

- ◆ $L = 8000$ bits (1000 bytes)
- ◆ Bandwidth $R = 2$ Mbps
- ◆ *If ignore propagation delay:* i.e. when last bit of a packet left A, it arrives at router-1 instantly
delay = $3 \times L/R = 12$ msec

Packet-switching: store-and-forward

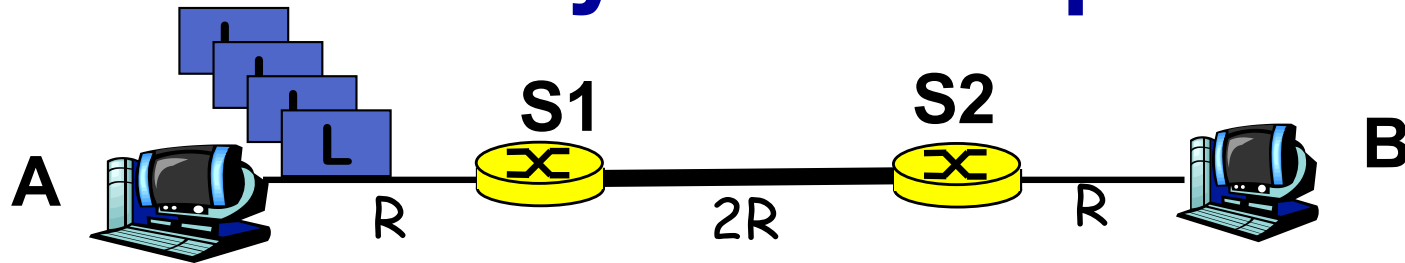


Example 2:

- ◆ A sends 5 packets to B
- ◆ $L = 8000$ bits, $R = 2$ Mbps
 - *Ignore propagation delay*
- ◆ How long does it take starting from A sending the first bit of first packet till B receives the last bit of the last packet?

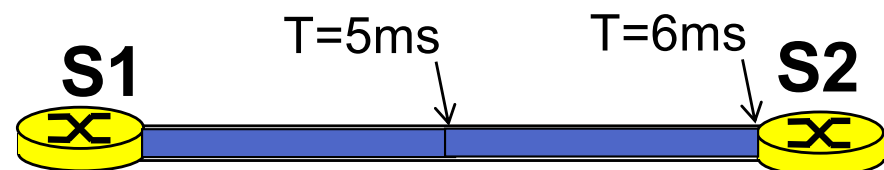
What if one takes into account the propagation delay?

Lets try an example

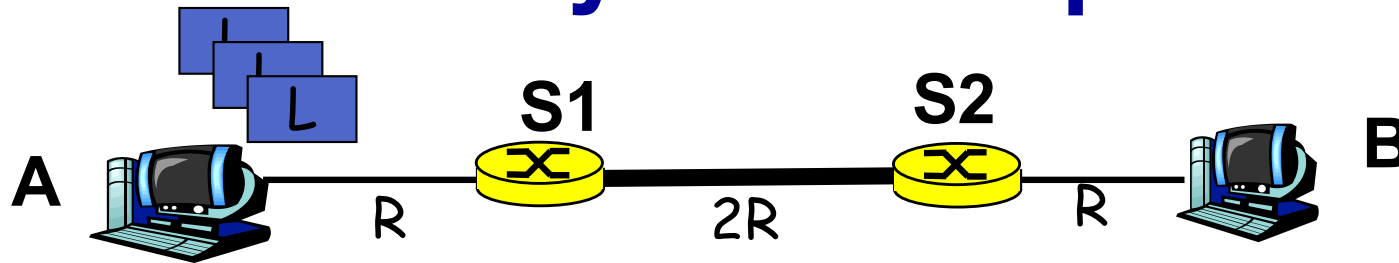


- ◆ $L = 4000$ bits, $R = 2$ Mbps, A sends 4 packets
 - $L/R = 2\text{msec}$, $L/2R = 1\text{ms}$
- ◆ Propagation delay: 2msec for each link
- ◆ When will first packet get to S2?
 - $A \Rightarrow S1: L/R + D_{\text{propagation}} = 2\text{ms} + 2\text{ms}$
 - $S1 \Rightarrow S2: L/2R + D_{\text{propagation}} = 1\text{ms} + 2\text{ms}$

When 1st packet arrives at S2, where is the 2nd packet?



Lets try an example

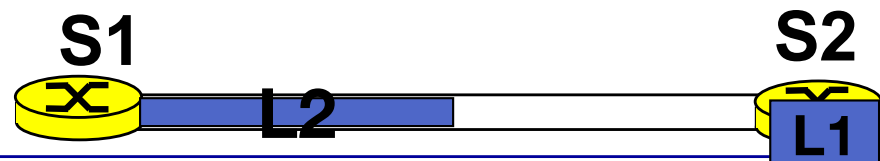


- ◆ $L/R = 2\text{msec}$, $L/2R = 1\text{ms}$
- ◆ Propagation delay: 2msec for each link
- ◆ When 1st packet arrives at S2, where exactly is the 2nd one?
→ $T=7\text{ms}$

$T=4\text{ms}$: last bit leaves A

$T=6\text{ms}$: last bit arrives S1

$T=7\text{ms}$: last bit leaves S1



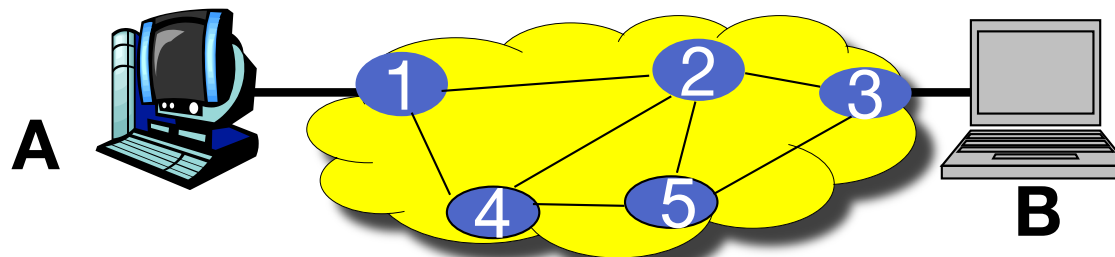
Network latency

- ◆ The time to send **1** packet from host A to B
 - sum of delays across each hop along the path

$$Delay_{A-B} = Delay_{A-1} + Delay_{1-2} + Delay_{2-3} + Delay_{3-B}$$

- ◆ **RTT**: round-trip-time

$$RTT_{AB} = Delay_{A-B} + Delay_{B-A}$$



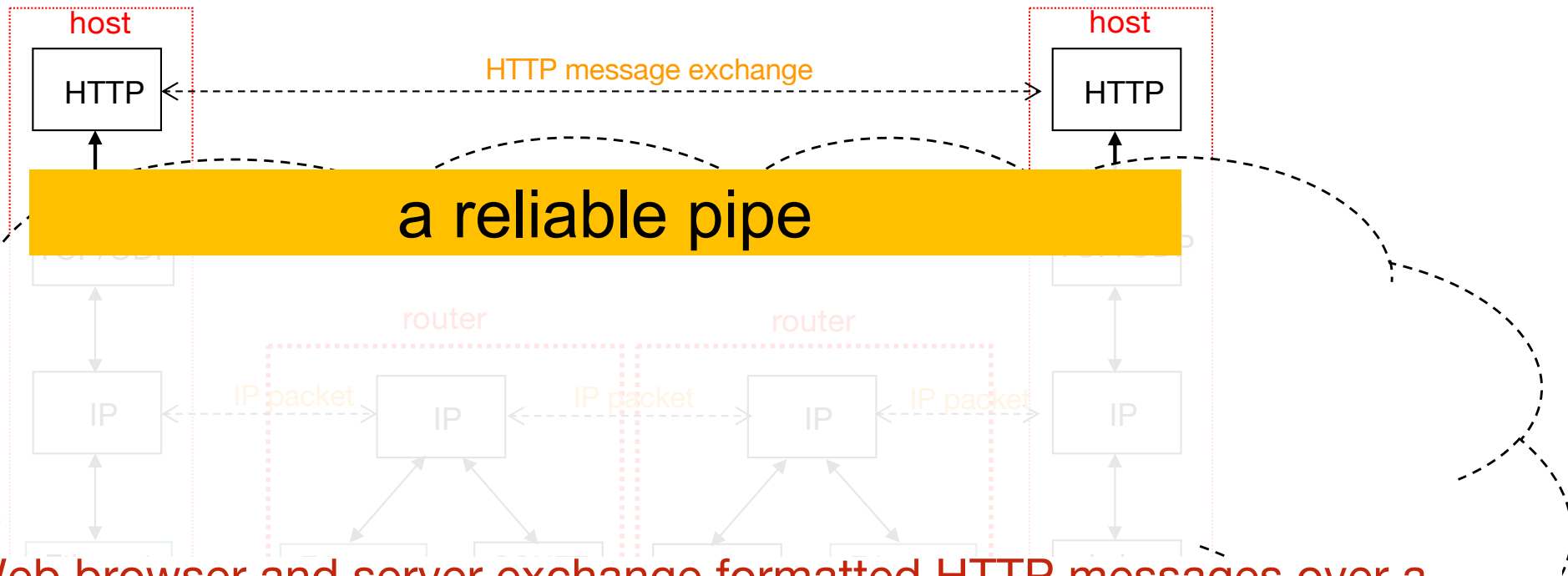
What we covered today

- ◆ Internet: made of a huge number of hosts, routers, wired and wireless links
- ◆ Hosts: run application protocols to exchange data packets with each other
- ◆ Routers: run bunch of protocols to move all packets towards their destinations
- ◆ Why protocols are layered
- ◆ How to calculate packet delays as they move across a packet-switched network

Lecture 1 Review: layered protocol architecture

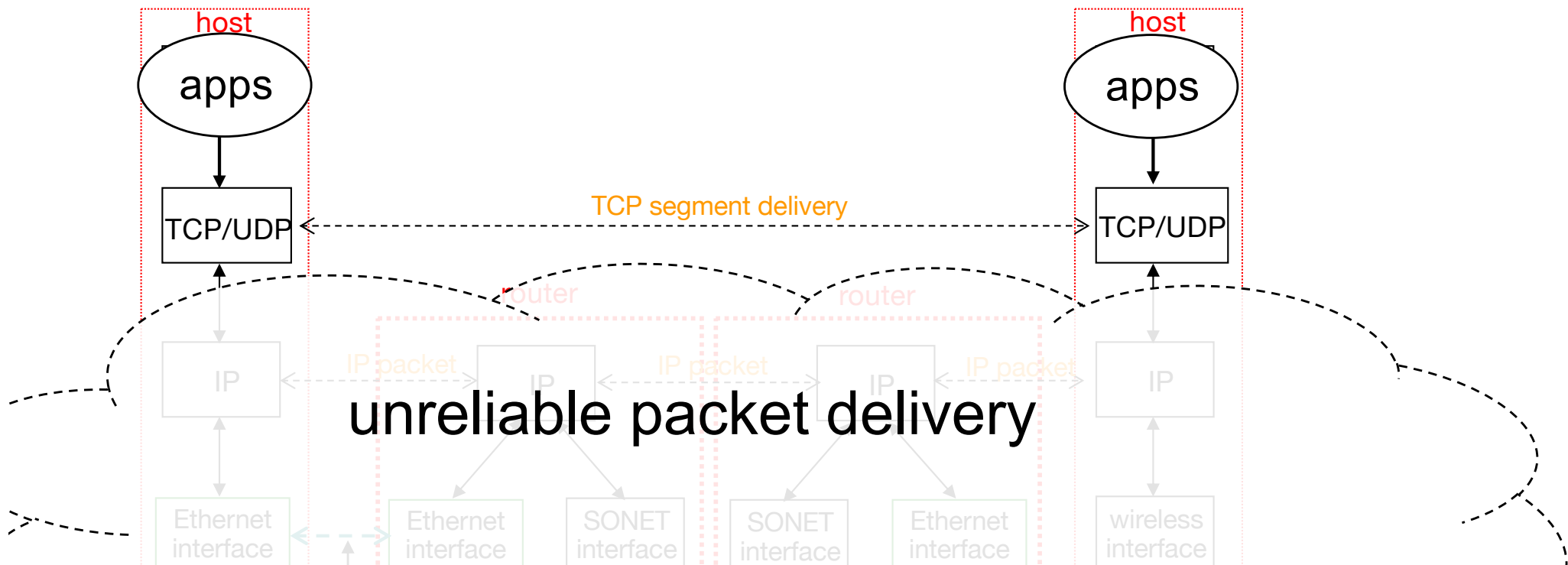
- ◆ Concepts:
 - **Internet:** made of a huge number of hosts and routers, interconnected by physical and wireless links
 - **Host:** a computer running applications and bunch of protocols to let apps exchange data with each other
 - **Router:** a packet switch running bunch of protocols to move packets toward their destinations
- ◆ Protocols are organized in layers:
 - Application protocols
 - Transport protocols
 - Network protocols
 - Link layer protocols
 - Physical layer
- ◆ How to calculate packet delays as they move across one hop

Application protocol's view of the world



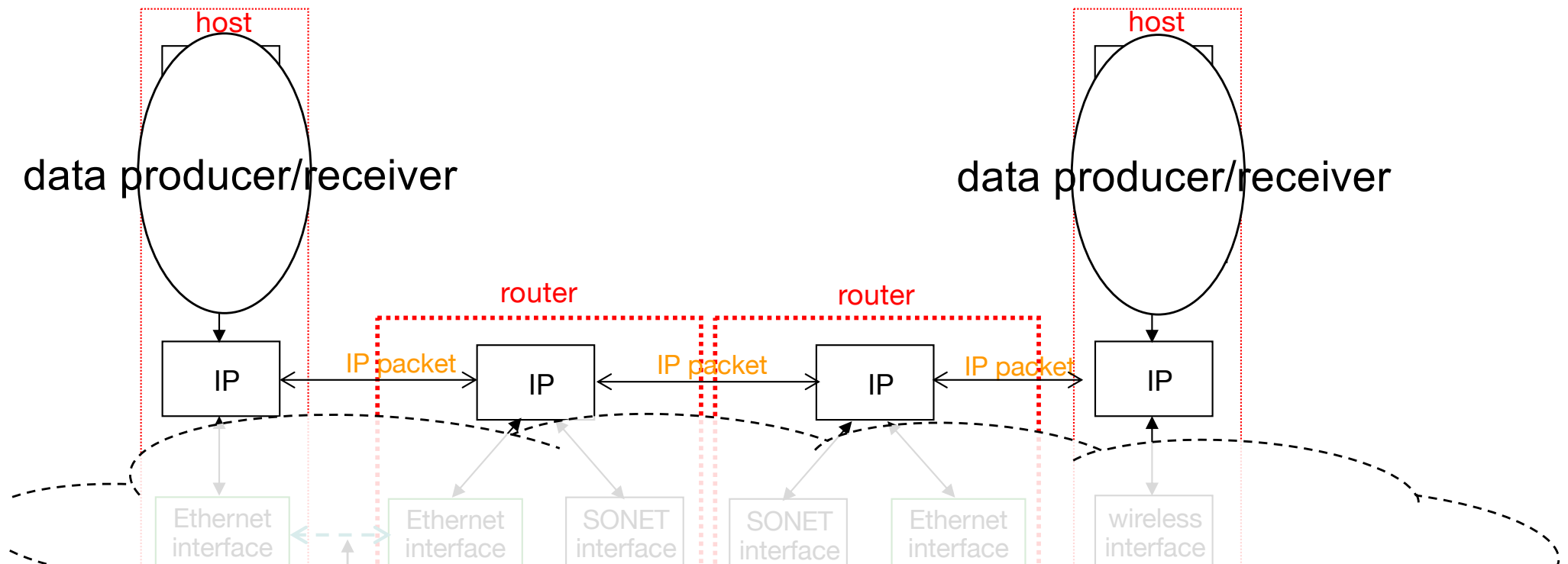
- ◆ Web browser and server exchange formatted HTTP messages over a reliable pipe
- ◆ As an application protocol, HTTP only concerns with the message's presentation format
- ◆ Application decides where msgs should be delivered to
 - The receiving end is identified by its name, which gets translated to IP address

Transport protocol's view of the world



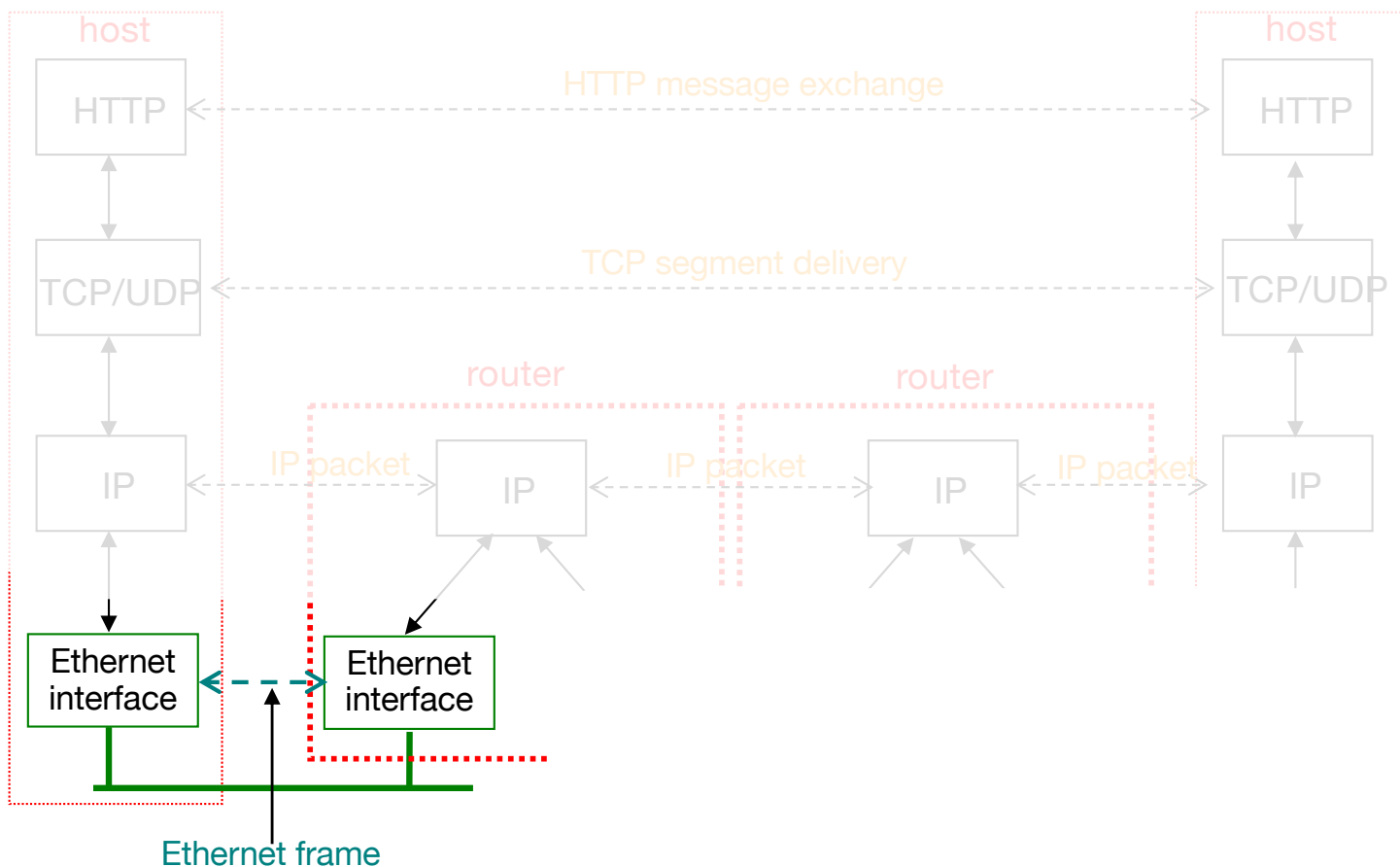
- ◆ A transport protocol receives data blobs from an application process, delivers them to the destination process (reliably)
 - Dest. Process is identified by IP address + (trans)port number
- ◆ It runs between two processes over an unreliable network (where packets can be garbled, lost, or reordered)

Network protocol's view of the world



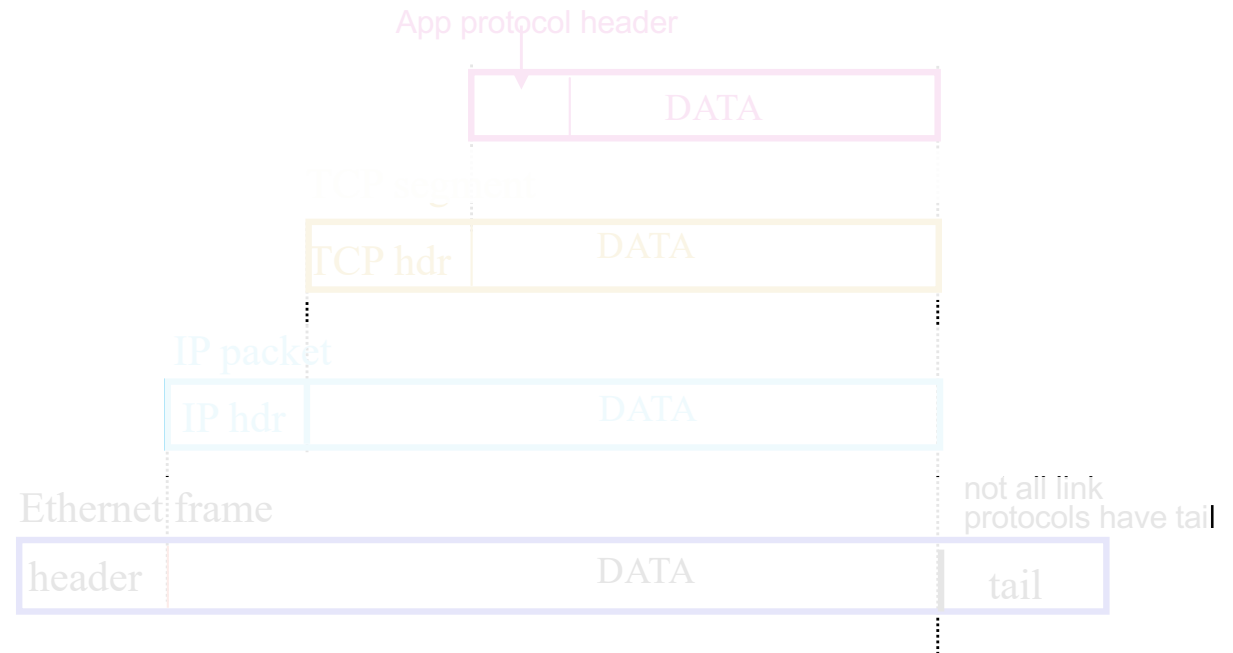
- ◆ Network protocol, IP, sees all IP-speaking nodes
- ◆ It receives data segments, delivers each of them to its destination IP address (with its best effort)
 - A router forwards packets, without looking inside IP envelope

Link layer protocol's view of the world



- ◆ A link layer protocol delivers data frames between two physically connected nodes
 - ◆ A link-layer header is added at sending node, removed by the receiving node
 - ◆ When a packet moves through the network across multiple hops: link-layer header is added and removed multiple times

Layered protocol implementation

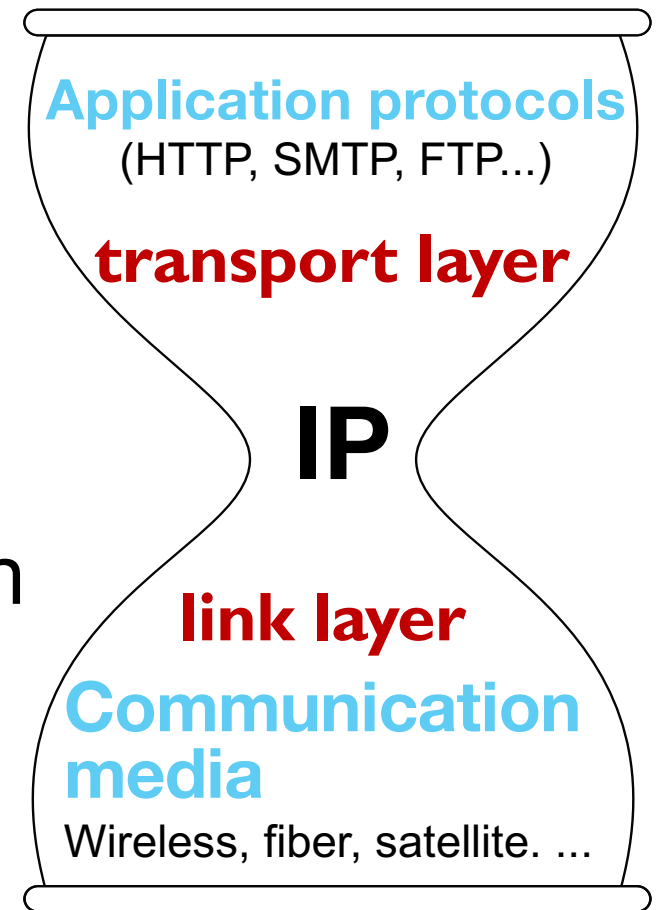


- ◆ protocol header: contains the information one writes on the “envelope”
- ◆ all the information, and *only* the information, that’s needed to carry out the protocol’s functionality

One more question: why 5 layers?

- ◆ Two layers are taken as given
 - Multiple different **application protocols**
 - Multiple different **physical communication media types**
- ◆ **IP**: the span layer
 - Connecting up all nodes
- ◆ **Link layer**: adaptation between IP and physical media
- ◆ **Transport**: adaptation between what apps want and what IP offers

5-year protocol stack

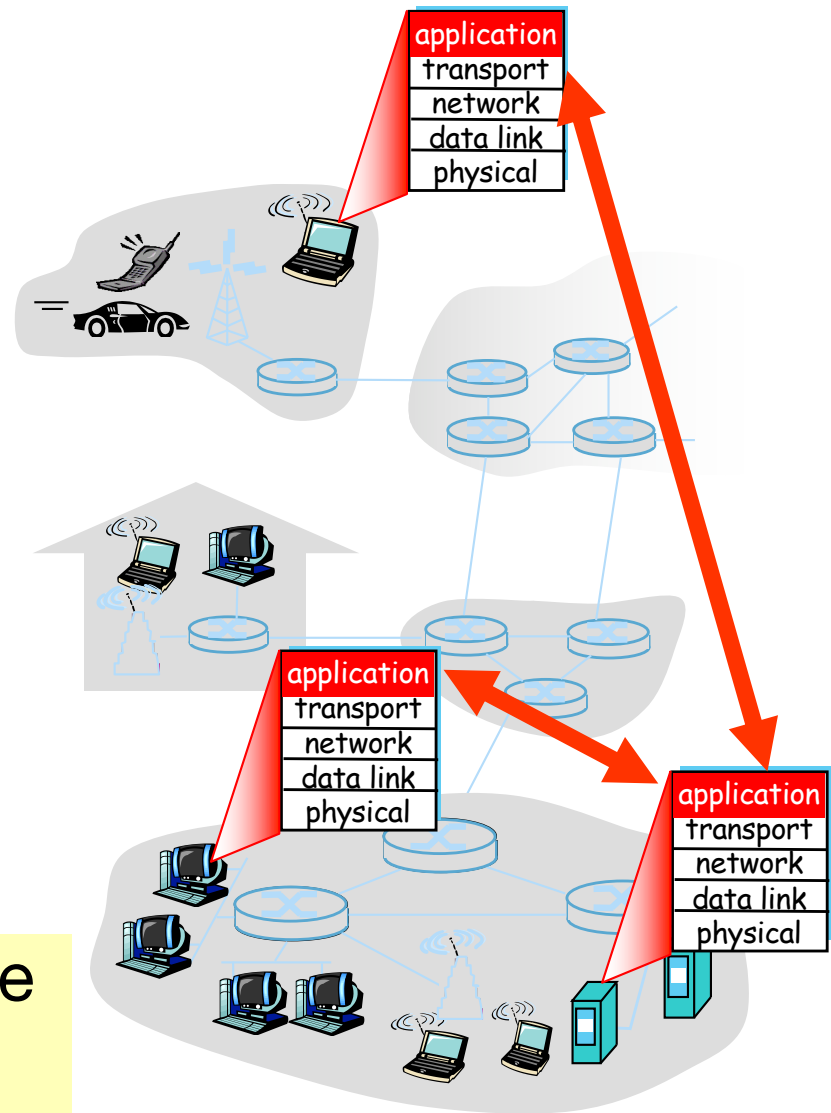


**Network applications:
how different parties reach each other**

Some popular network applications

- ◆ Email
- ◆ Web
- ◆ WhatsApp
- ◆ BitTorrent (P2P file sharing)
- ◆ Online Games
- ◆ YouTube
- ◆ Virtual Conferencing

Application processes communicate with each other using application protocols



Client-server application communication model

servers:

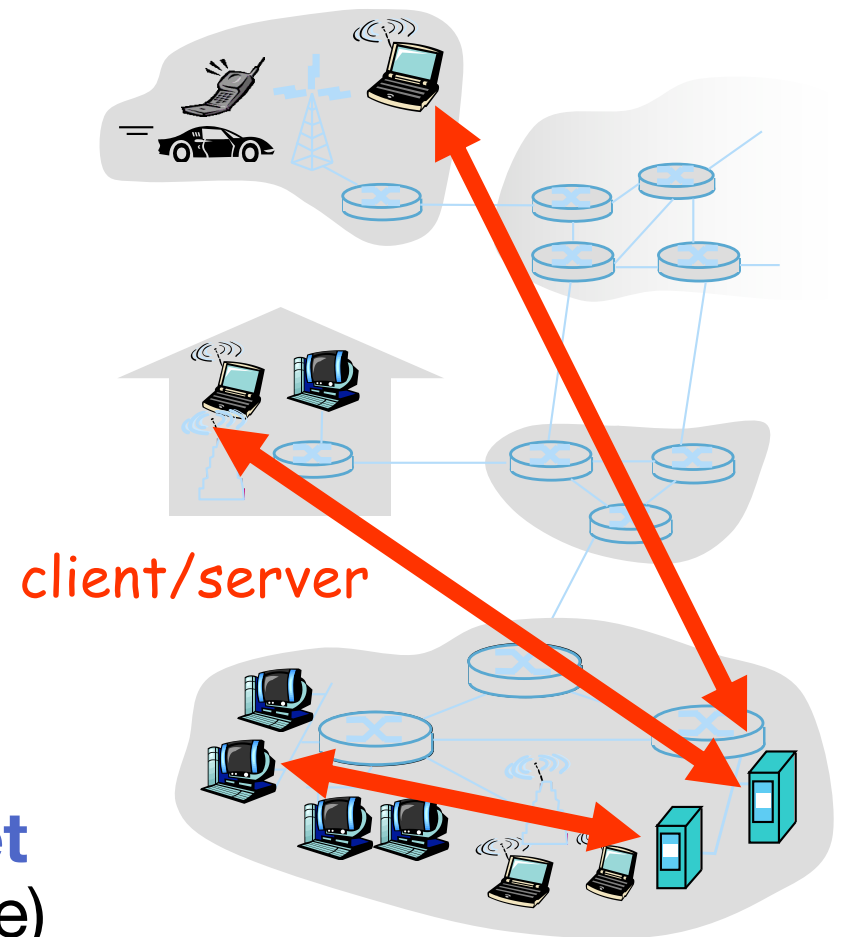
- ◆ Reachable by IP address
- ◆ **always-on**, waiting for incoming requests from clients

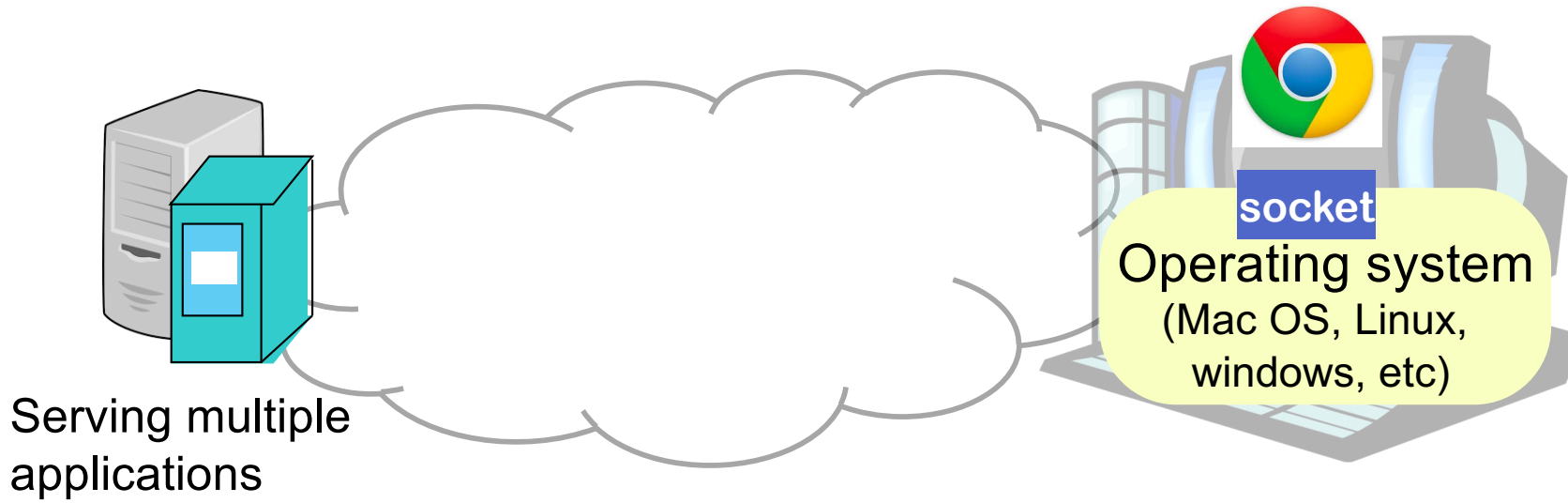
clients:

- ◆ Initiate communication with server

Q: How does a client process *identify* the server process with which it wants to communicate?

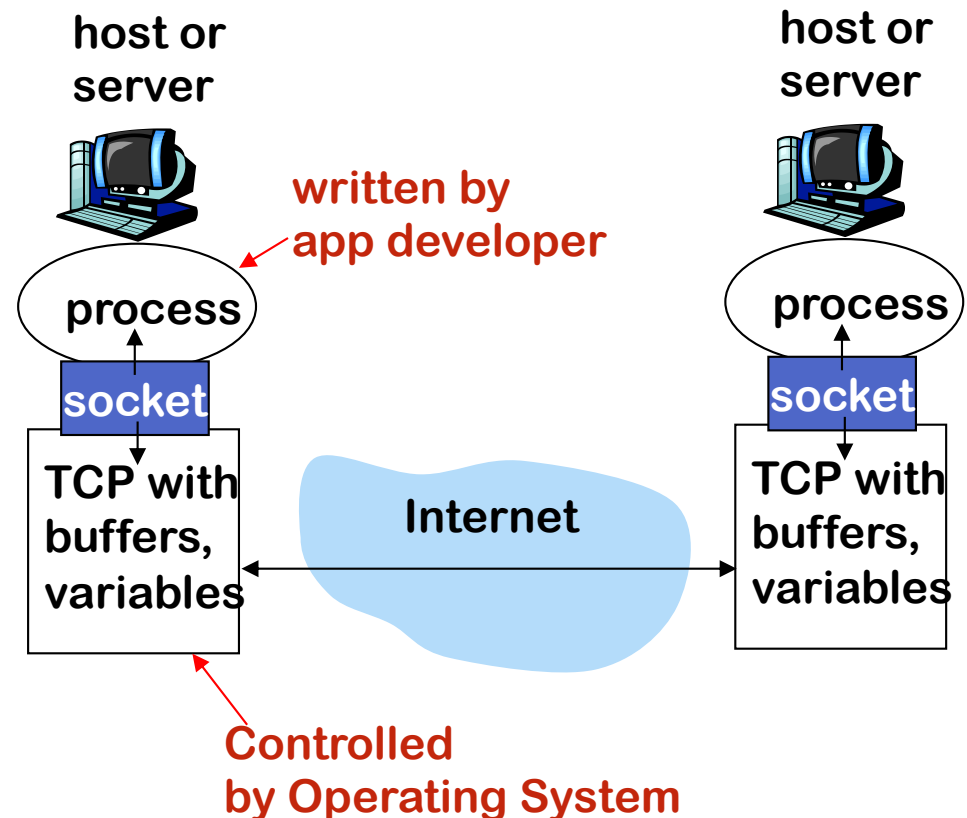
A: Using port numbers via the **socket API (Application Program Interface)**





Socket

- ◆ **Process**: program running on a host
- ◆ Between different hosts: Processes communicate through an **application-layer protocol**
- ◆ A process sends/ receives messages to/from its **socket**
- ◆ A socket analogous to a door:
 - sending process shoves message out of the door
 - transport protocol brings message up to the socket at receiving process



What is "socket"

- ◆ A set of system function calls

socket (): Create a socket

bind(): bind a socket to a local IP address and port #

connect(): initiating connection to another socket

listen(): passively waiting for connections

accept(): accept a new connection

Write(): write data to a socket

Read(): read data from a socket

Close()

host or
server

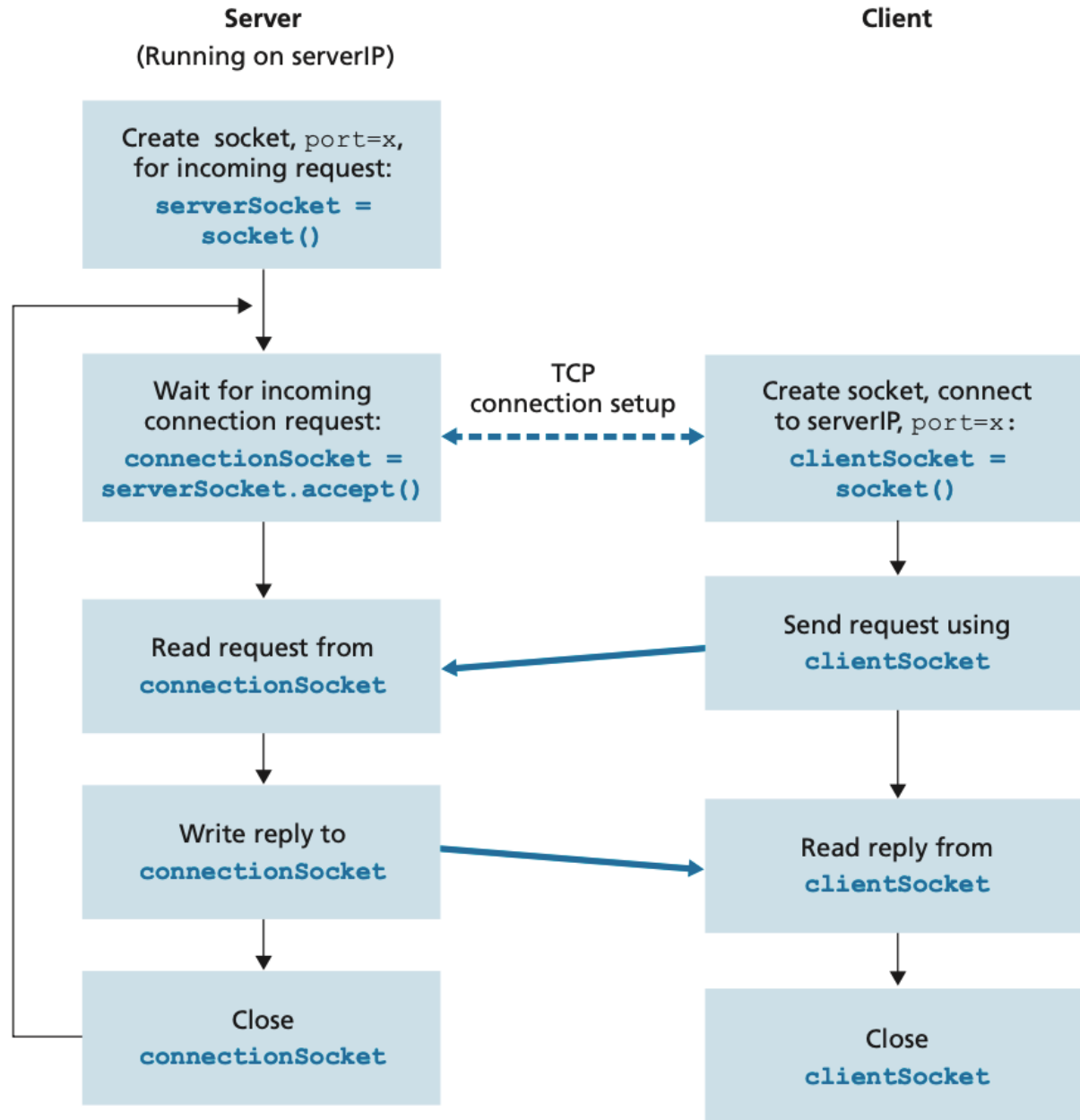


process

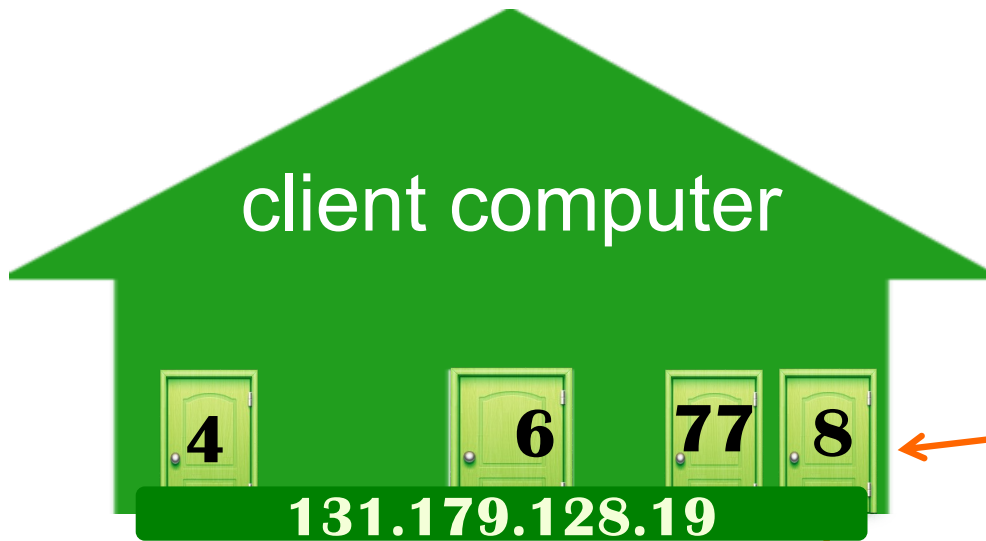
socket

TCP with
buffers,
variables

What is "socket"



Socket: analogous to a door



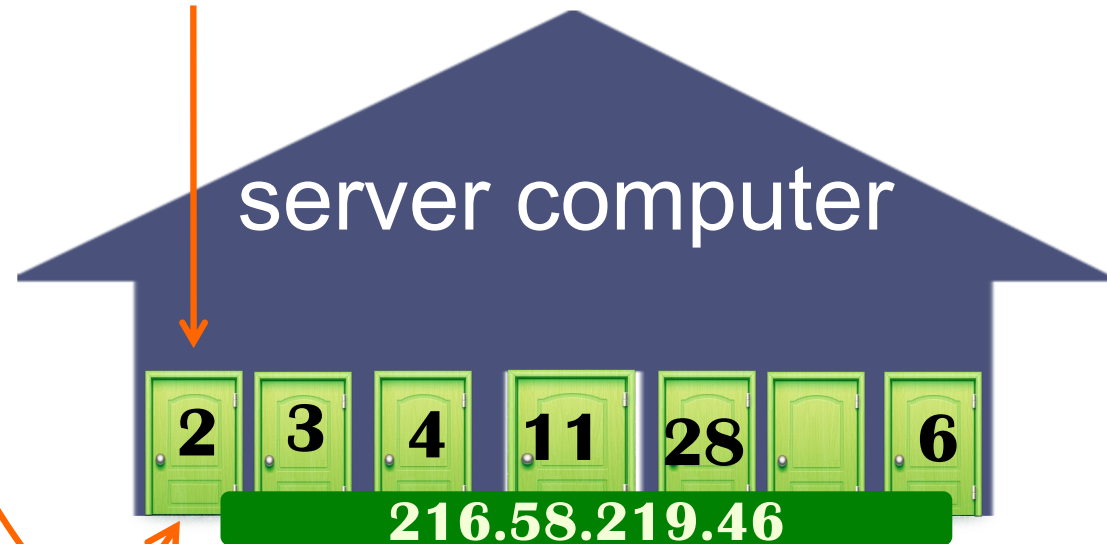
connect()

socket(): create a “door”
bind(): tie the door to a
[local IP addr, port#] pair

close(): delete/remove a door

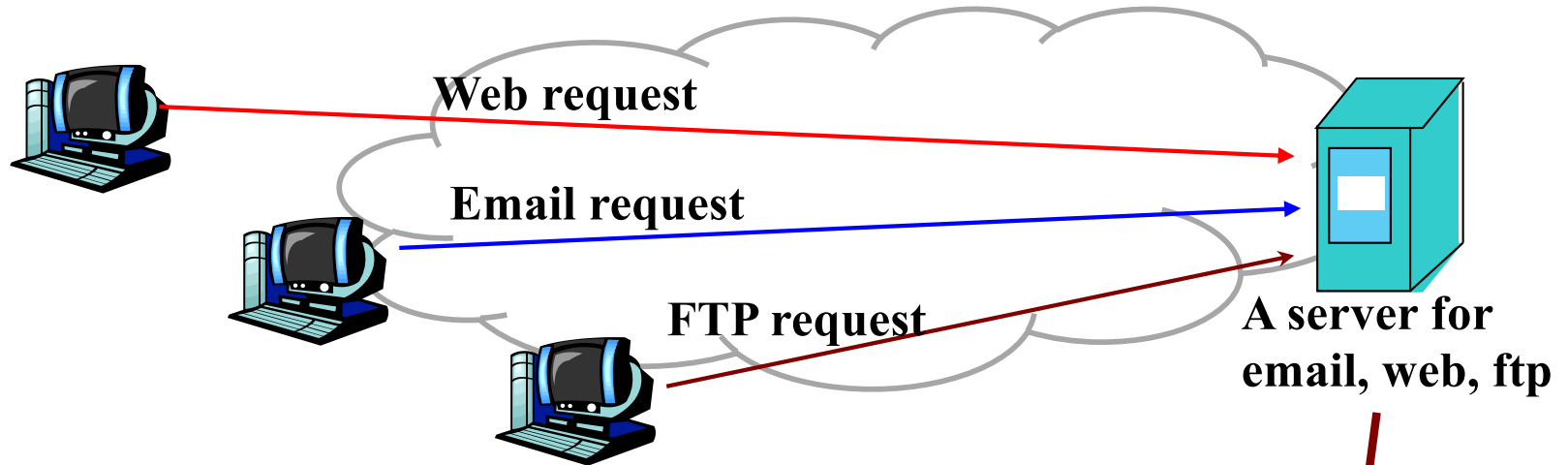
A couple other questions:

- What is the server’s IP address?
 - Website name → IP address
- What port number to use?
 - Client: assigned by OS
 - Server: defined standards



listen(): start waiting for incoming
packet with matching destination port#

A quick comment about "port"



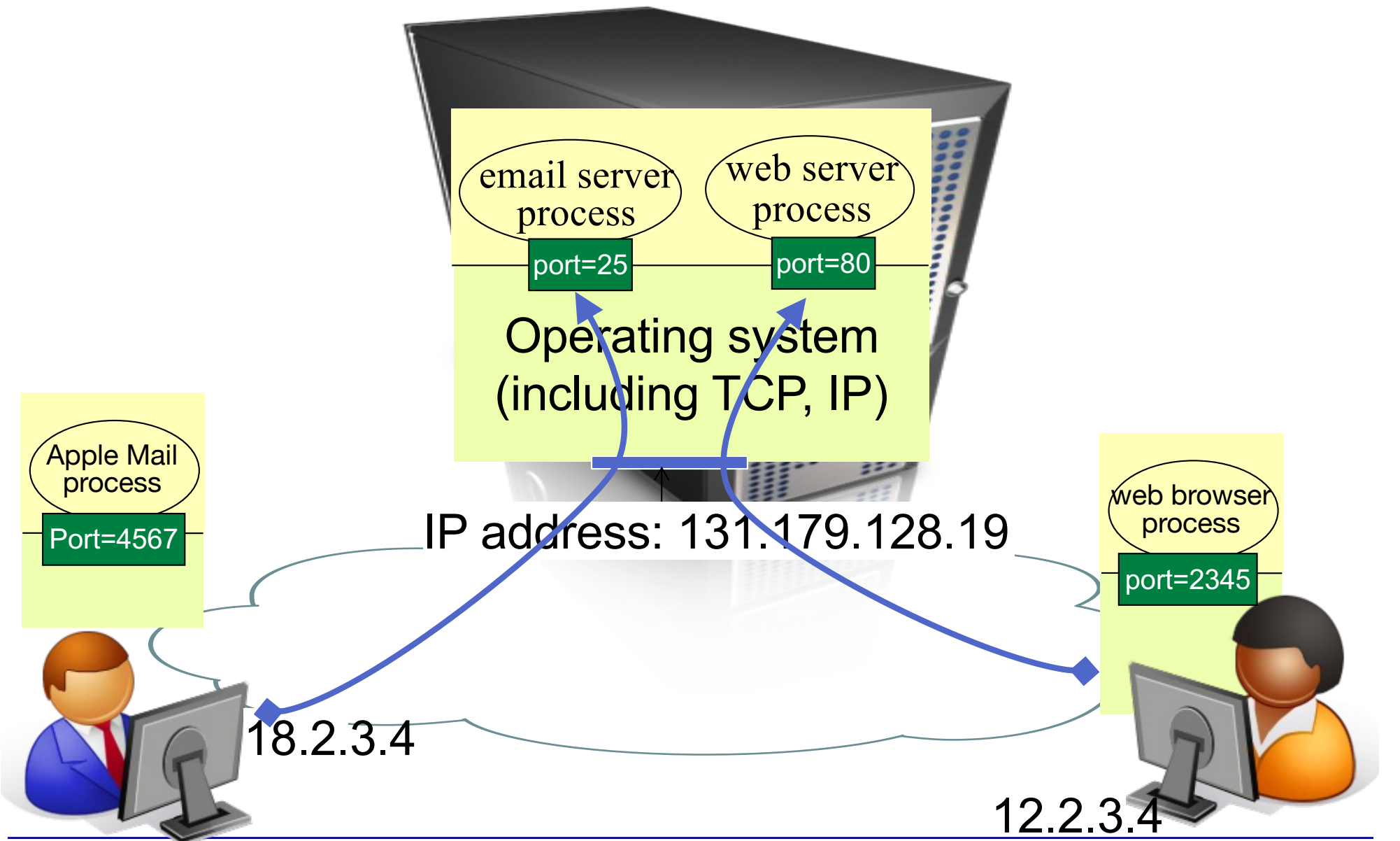
- ◆ Web, Email, FTP all connect client to server by TCP
- ◆ The server tells which client wants which service by *well-defined port number*:
 - Web: port 80, ftp 21, mail 25

ftp:waiting for packets to port21

email process: port25

Web server: port80

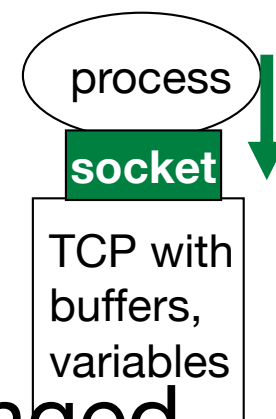
IP address, TCP connection, port number, processes, and sockets



Applications

So far we've talked

- ◆ Application process (executing application program)
- ◆ Application protocol (used by application processes to exchange data)
- ◆ Exactly how data is exchanged
 - Socket
 - Transport protocol
- ◆ Lets look at exactly what data is exchanged



Web and HTTP

- ◆ **Web page**: normally consists of
 - **base HTML-file**, which includes
 - several referenced **objects**
- ◆ An object can be another HTML file, JPEG image, Java applet, audio file,...
- ◆ Each object is addressable by a **URL** (Universal Resource Locator)

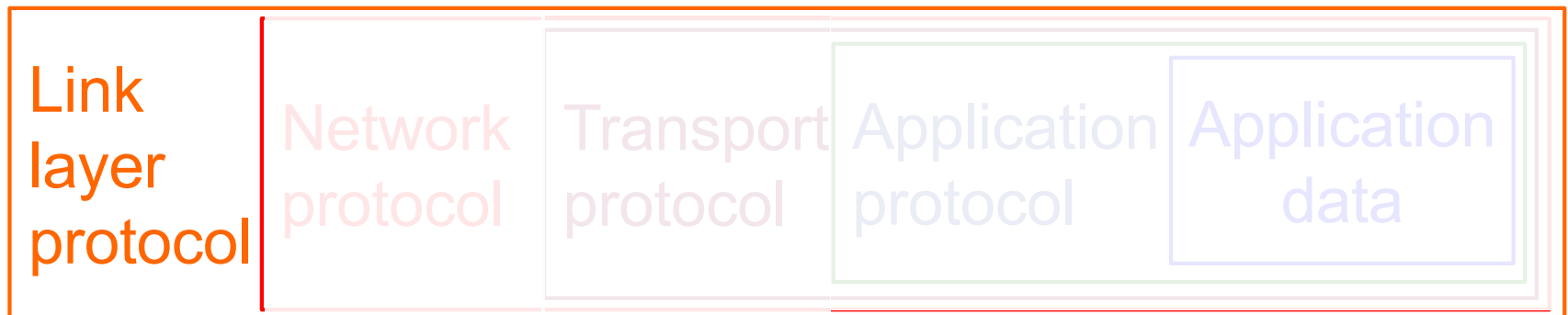
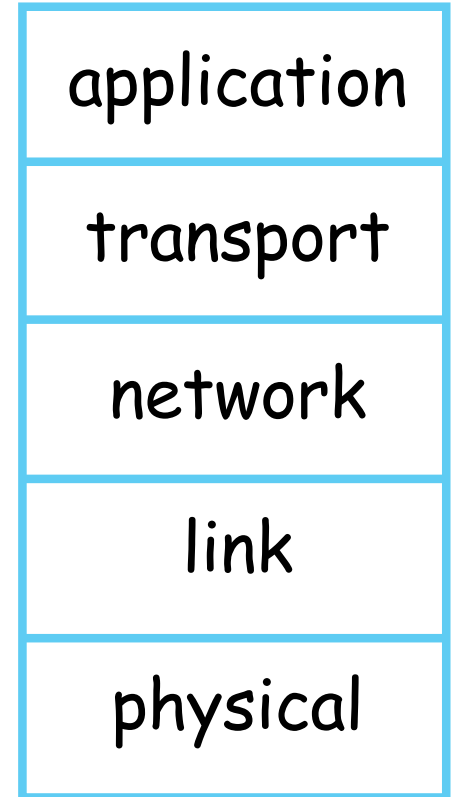
`http://www.someschool.edu:port#/someDept/pic.gif`

↑
Application protocol

host name

path name

What protocol “layer” really means



Acknowledgment

- ◆ Slides adapted from S24 CS118 instructed by Prof. Lixia Zhang